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Detecting the Change of Clustering Structure in Categorical Data Streams

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Abstract

Analyzing clustering structures in data streams can provide critical information for making decision in real-time. Most research has been focused on clustering algorithms for data streams. We argue that, more importantly, we need to monitor the change of clustering structure online. In this paper, we present a framework for detecting the change of critical clustering structure in *categorical data streams*, which is indicated by the change of *the best number of clusters* (Best K) in the data stream. The framework extends the work on determining the best K for static datasets (the BkPlot method) to categorical data streams with the help of a *Hierarchical Entropy Tree* structure (HE-Tree). HE-Tree can efficiently capture the entropy property of the categorical data streams and allow us to draw precise clustering information from the data stream for high-quality BkPLots. The experiments show that with the combination of HE-Tree and the BkPlot method we are able to efficiently and precisely detect the change of critical clustering structure in categorical data streams.

1 Introduction

With the deployment of wide-area sensor systems and Internet-based continuous-query applications, processing stream data has become a critical task. As an important method in data analysis, recently clustering has attracted more and more attention in analyzing and monitoring streaming data [19, 2]. The initial research has shown that clustering stream data can provide important clues about the new emerging data patterns so that the decision makers can predict the coming events and react in near real time. Stream data clustering is especially important to the time-critical areas such as disaster monitoring, anti-terrorism, and network intrusion detection. As many of such applications also include a large amount of categorical data, clustering the categorical data streams becomes an interesting and challenging problem. Surprisingly, very few [4] have addressed the problems related to clustering categorical data streams.

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Most research has been focused on how to design online clustering algorithms for data streams. However, cluster analysis includes not only the clustering phase, but also the cluster evaluation and validation [24], which determines the critical clustering structure, such as the best K number of clusters. For stream clustering, most algorithms like Coolcat [4] assume that the best K number of clusters is given, which inappropriately simplifies the problem. One of the primary goals to clustering the data streams is to monitor the change of clustering structure may indicate possible new events. Therefore, distinguishing the difference between the clustering structures in different stages of data streams becomes critical, yet very challenging, to data stream applications. Surprisingly, none has addressed how to online evaluate the cluster structure, particularly, the best K clusters for data streams.

The change of critical clustering structure in data streams involves three aspects: the drifting of the cluster center caused by the increasing size of cluster, new emerging clusters, and disappearing clusters caused by the convergence of growing clusters. We observed that the the latter two aspects can be indicated by the "Best K" number of clusters [11]. Monitoring the change of critical clustering structure is important since it may correlate to some important events in the applications. For example, a network attack may correlate to the change of clustering structure in data streams. In this context, we are more interested in detecting both the emerging of clusters and the disappearing of clusters. The former may indicate some new events (the attacks) are going to happen, and the later implies two possible situations: 1) some clusters grow big and become merged, which may indicate that some correlated event may become significant and will occur more and more frequently; 2) the previously identified clusters are actually part of outliers. Fast detection of the change of critical clustering structure can eliminate most costly cluster evaluation work and allow to efficiently monitor the data streams− we only analyze individual clusters when the change of clustering structure is detected. E.g., we can visualize the new emerging clusters or the converged clusters, which requires more computational

resource and human interaction in practice.

In this paper, we will focus on detecting the change of clustering structure for *categorical data streams*. Extending the previous work on finding the best K for static categorical datasets (the BkPlot method [11]), we propose a method for detecting the change of best K clusters in categorical data streams. The original BkPlot method is an optimal method for identifying the candidate best Ks for a given dataset. Due to the NP-hard complexity in generating optimal BkPlots, only approximate BkPlots can be applied in practice. The ACE algorithm is proposed in paper [11] to generate high-quality approximate BkPlots. However, due to its still high complexity $O(N^2)$, it is impossible to apply the algorithm directly on data streams.

The key idea is based on the design of a summarization tree structure, called Hierarchical Entropy Tree (HE-Tree for short). HE-Tree utilizes a small amount of memory to summarize the entropy property of the data streams, and groups the data records into a bunch of sub-clusters located at the HE-Tree leaf nodes. The extended ACE algorithm is able to handle the snapshot sub-clusters (often a few hundreds) and generate an approximate snapshot BkPlot for identifying the Best K at certain time interval. The difference between the clustering structures can be conveniently identified by comparing the distinctive points on the snapshot BkPlots.

The rest of the paper is organized as follows. Section 2 sets down the notations and gives the concepts in entropy-based categorical clustering. Section 3 briefly introduce the BkPlot for finding the best K in static categorical datasets. In section 4, we develop the HE-Tree structure and describe its working mechanism. In section 5, we propose the framework of detecting the change of clustering structure in categorical data streams. The experimental results are shown in section 6. Finally, we review the related work of categorical clustering and stream clustering, and conclude our work.

2 Entropy-based Categorical Clustering

Clustering techniques for categorical data are very different from those for numerical data, mainly because of the definition of similarity measure. Most numerical clustering techniques have been using distance functions, for example, Euclidean distance, to define the similarity measure. However, there is no such inherent distance meaning between the categorical values.

In contrast to the distance-based similarity measure for pairs of data records, similarity measures based on the "purity" of a bulk of records seem more intuitive for categorical data. Entropy [13] is a well defined measure for the purity of dataset. Originally from information theory, entropy has been applied in various areas, such as pattern discovery [7], numerical clustering [12] and information retrieval [28]. Due to the lack of intuitive distance definition for categorical values, recently entropy has been applied in clustering categorical data [4, 25, 9, 14]. The initial results have shown that the entropy criterion can be very effective in clustering categorical data. Paper [25] also proves that the entropy criterion can be formally derived in the framework of probabilistic clustering models, which strongly supports that the entropy criterion is a meaningful and reliable similarity measure, particularly good for categorical data.

In entropy-based categorical clustering, the quality of clustering is essentially evaluated by the entropy criterion, namely, the *Expected Entropy* of clusters [4, 25]. Other variants, such as Minimum Description Length (MDL) [9] or mutual information [14, 3], turn out to be equivalent to the entropy criterion, as the paper [25] shows. We categorize all these approaches as entropy-based categorical clustering. The main goal of the entropy-based clustering algorithms is to find a partition that minimizes the expected entropy for K clusters. However, minimizing expected entropy is a NP-hard problem, thus it is computationally intractable even for a median-sized dataset. A common approach to solving this problem is approximation. Typically, in approximation algorithms we have to sacrifice some optimality to obtain the efficiency. Categorical data streams make it even harder to balance the two conflicting factors. With

Below, we first give the notations and definitions used in this paper, and then describe an important metric *Incremental Entropy*, which is used in HE-Tree construction and extended ACE algorithm.

2.1 Notations and Definitions Consider that a dataset $\mathbb{S}$ with N records and d columns, is a sample set of the discrete random vector $X = (x_1, x_2, \ldots, x_d)$. For each component x_j, $1 \leqslant j \leqslant d$, x_j takes a value from the domain A_j. A_j is conceptually different from $A_k (k \neq j)$. There are a finite number of distinct categorical values in $domain(A_j)$ and we denote the number of distinct values as $|A_j|$. Let $p(x_j = v)$, $v \in A_j$, represent the probability of $x_j = v$, we have the classical entropy definition [13] as follows.

$$\begin{aligned} H(X) &= \sum_{j=1}^{d} H(x_j) \\ &= -\sum_{j=1}^{d} \sum_{v \in A_j} p(x_j = v) \log_2 p(x_j = v) \end{aligned}$$

Since $H(X)$ is estimated with the sample set $\mathbb{S}$, we define the estimated entropy as $\hat{H}(X) = H(X|\mathbb{S})$, i.e.

$$\begin{aligned}\hat{H}(X) &= H(X|\mathbb{S}) \\ &= -\sum_{j=1}^{d}\sum_{v\in A_j} p(x_j = v|\mathbb{S}) \log_2 p(x_j = v|\mathbb{S})\end{aligned}$$

Suppose the dataset $\mathbb{S}$ is partitioned into K clusters. Let $C^K = \{C_1, \ldots, C_K\}$ represent a partition, where C_k is a cluster and n_k represent the number of records in C_k. The classical entropy-based clustering criterion tries to find the optimal partition, C^K, which maximizes the following entropy criterion [6, 8, 25].

$$O(C^K) = \frac{1}{d}\left(\hat{H}(X) - \frac{1}{N}\sum_{k=1}^{K} n_k \hat{H}(C_k)\right)$$

Since $\hat{H}(X)$ is fixed for a given dataset $\mathbb{S}$, maximizing $O(C^K)$ is equivalent to minimize the item $\frac{1}{dN}\sum_{k=1}^{K} n_k \hat{H}(C_k)$, which is named as the "expected entropy" of partition C^K. Let us notate it as $\bar{H}(C^K)$. For convenience, we also name the varying part $n_k\hat{H}(C_k)$ as the "Weighted Entropy" of cluster C_k.

Entropy criterion is especially good for categorical clustering due to the lack of intuitive distance function for categorical values. While entropy criterion can also be applied to numerical data [12] if the numerical data is appropriately discretized, it is unable to describe most of the inherent geometric clustering features for the numerical data.

2.2 Incremental Entropy While expected-entropy describes the average intra-cluster quality, incremental entropy is a measure used to describe the similarity between any two clusters. We begin with the observation about the change of expected entropy when merging two clusters. Intuitively, merging the two clusters that are similar in the inherent structure will not increase the disorderliness (expected-entropy) of the partition, while merging dissimilar ones will inevitably bring larger disorderliness. Therefore, this increase of expected entropy has some correlation with the similarity between clusters. It is, thus, necessary to formally explore the property of merging clusters. Let $C_p \cup C_q$ represent the mergence of two clusters C_p and C_q, and C_p and C_q have n_p and n_q members, respectively. Suppose that the number of clusters is $K+1$ before the mergence happens. By the definition of expected entropy, the difference between $N\hat{H}(K)$ $N\hat{H}(K+1)$ is only the difference between the weighted entropies, $(n_p+n_q)\hat{H}(C_p \cup C_q)$ and $n_p\hat{H}(C_p) + n_q\hat{H}(C_q)$. Intuitively, since merging always increases the entropy, we have the following relation for the weighted entropies.

PROPOSITION 1. $(n_p + n_q)\hat{H}(C_p \cup C_q) \geqslant n_p\hat{H}(C_p) + n_q\hat{H}(C_q)$

The detailed proof about above proposition can be found in [11]. We name $I_m(C_p, C_q) = (n_p + n_q)\hat{H}(C_p \cup C_q) - (n_p\hat{H}(C_p) + n_q\hat{H}(C_q)) \geqslant 0$ as the "*Incremental Entropy* (IE)" of merging the clusters C_p and C_q. Note that $I_m(C_p, C_q) = 0$ suggests that the two clusters have the identical structure – for every categorical value v_i in any arbitrary attribute x_j, $1 \leqslant i \leqslant |A_j|$, $1 \leqslant j \leqslant d$, we have $p(x_j = v_i|C_p) = p(x_j = v_i|C_q)$. The larger I_m is, the more different the two clusters are. IE plays an important role in constructing a hierarchical clustering scheme, where merging with IE measure is equivalent to minimize the expected-entropy criterion. We also use IE as a major measure in HE-Tree operations.

3 BkPlot for Determining the "Best K" for Categorical Clustering

In order to help understand the entire framework for detecting the change of clustering structure for categorical data streams, we briefly describe the BkPlot method for determine the candidate best K for static datasets. For detailed description and analyze, please refer to paper [11].

Traditionally, statistical validity indices based on geometry and density distribution are applied in clustering numerical data [21]. A typical index curve consists of the statistical index values for different K number of clusters. The Ks at the peaks, valleys, or distinguished "knees" on the index curve, are regarded as the candidates of the optimal number of clusters (the best K). BkPlot method tries to find such kind of index for categorical data clustering.

If the neighboring partitions are defined as two clustering results having K and $K+1$ number of clusters, respectively, the basic idea of BkPlot is to investigate the entropy difference between any two optimal neighboring partitions. Let the expected-entropy of the optimal partition notated as $\bar{H}_{opt}(C^K) = \min\{\bar{H}_i(C^K)\}$, where i is the index of all possible K-cluster partitions. The curve of $\bar{H}_{opt}(C^K)$ is identified as a smoothly decreasing curve, i.e., $\bar{H}_{opt}(C^K) \geq \bar{H}_{opt}(C^L)$, for $K < L$, without any distinguished peaks, valley, or knees, from which we cannot effectively identify the best K.

However, the special meaning behind the entropy difference of the neighboring partitions enables us to explore the best K. Let the increasing rate of entropy between the optimal neighboring partitions defined as $I(K) = \bar{H}_{opt}(C^K) - \bar{H}_{opt}(C^{K+1})$. We identify that $I(K)$ implies two levels of difference between the neighboring partitions.

- $I(K)$ is the level of difference between the two

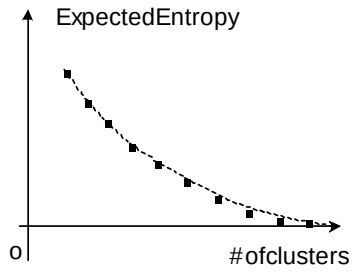

Figure 1: Sketch of expected entropy curve.

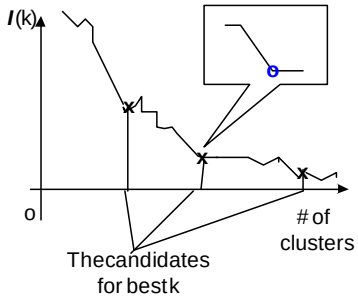

Figure 2: Sketch of ECG graph.

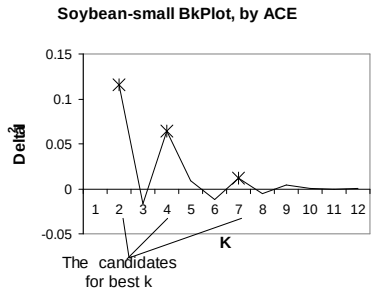

Figure 3: Finding the best k with BkPlot (for soybean-small dataset).

neighboring schemes. The larger the difference, the more significant the clustering structure is changed by reducing the number of clusters by 1.

- Consider $I(K)$ as the amount of impurity introduced from $K+1$-cluster scheme to K-cluster scheme. If $I(K) \approx I(K+1)$, i.e. K-cluster scheme introduces similar amount of impurity as $K+1$-cluster scheme does, the change of clustering structure follows the "similar pattern, thus we can also consider that there is no significant difference from $K+2$-cluster partition to K-cluster partition.

We define the differential of expected-entropy curve as "Entropy Characteristic Graph (ECG)" (Figure 2). An ECG shows that the similar partition schemes with different K are at the same "plateau". From plateau to plateau there are the critical points implying the significant change of clustering structure, which could be the candidates for the best Ks.

The common way to automatically identify such critical knees on ECG is to find the peaks/valleys at the second-order differential of ECG. Since an ECG consists of a set of discrete points, we define the second-order differential of ECG as $\delta^2 I(K) : \delta I(K) = I(K) - I(K+1)$ and $\delta^2 I(K) = \delta I(K-1) - \delta I(K)$ to make K aligned with the critical points. These critical points are highlighted in the second-order differential of ECG, which is named as "Best-K Plot (BkPlot)".

Exact BkPlots cannot be achieved in practice, since $I(K)$ is based on the optimal K-cluster scheme which involves entropy minimization. However, since we only pay attention to the peak/valley points, approximate but accurate BkPlots are possible to acquire. A hierarchical clustering algorithm ACE in [11] is proposed to generate such BkPlots, and we have shown in experiments that ACE is a robust method to generating the high-quality BkPlots. ACE also has a nice property that we only need to observe the peaks in the BkPlots generated by ACE to determine the best Ks. For interested readers, please find more details in [11].

ACE is initially designed for static datasets and the $O(N^2)$ complexity prevents it working directly on large datasets or data streams. The extended discussion [] also shows that ACE can run on samples of large datasets so that the generated sample BkPlots are consistent with the original one. However, sampling approach does not apply for the data stream scenario since the clustering structure is changing over time. In the next section, we will design a data summarization structure with the help of Incremental Entropy, the result of which can be combined with the extended ACE algorithm to generate high-quality BkPlots for data streams.

4 HE-Tree: Capture the Entropy Characteristics of Categorical Data Stream

In this section, we design the summarization structure – Hierarchical Entropy Tree (HE-Tree). The basic idea of HE-Tree is to coarsely but rapidly assign the records from the data stream onto hundreds of subclusters. Observing these subclusters will give us a precise estimation of the clustering structure. HE-Tree determines the subclusters only based on the previously processed data records, and the clustering structure of the subclusters can automatically adapt to the new coming data records. The criterion for forming a subcluster is minimizing the expected-entropy of the subclusters – certainly, this minimization is only locally optimal, but it generates good global approximations. To fast locate a subcluster for a coming new record, we organize the subclusters in a tree, i.e., HE-Tree.

A HE-Tree consists of two key components:

1. *HE-node structure*, which summarizes the entropy characteristics of a group of records and facilitate fast processing of stream data items;

2. *Incremental-Entropy based lookup/assigning algorithm*, which helps to adapt the changing clustering structure.

Given the fixed height h and fanout f, HE-Tree is constructed in two stages:

1. *growing stage*, which happens at the beginning of processing data stream;

2. *absorbing stage*, which absorb the new coming items to the subclusters at the leaf nodes, when the tree is full.

We first give the structure of the HE-Tree node, which includes the structures for fast entropy calculation. After that, we will focus on the construction algorithms of the HE-Tree.

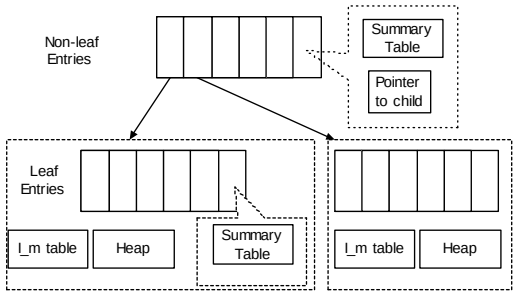

Figure 4: The structure of HE-Tree.

4.1 Structure of HE-Tree Summary Table. Summary table is used to maintain the fast calculation of the entropy $\hat{H}(C_k)$ and each node in the HE-Tree maintains one summary table Since computing cluster entropy is based on counting the occurrences of categorical values in each column, summary table is used to keep the counters for each cluster. For each categorical value $v_{ij} \in A_j$, we have an element $T[v_{ij}]$ in summary table as its corresponding counter. Therefore, if the average column cardinality is m, a summary table keeps dm counters. Obviously, summary table also has the following property.

PROPOSITION 2. *When two clusters are merged, the sum of the two summary tables becomes the summary table for the new cluster.*

Nodes in HE-Tree. HE-Tree is a balanced tree similar to B-tree, where each node has f entries and the entries in the leaf nodes represents the n_c subclusters. As shown in Figure 4, each entry in leaf node contains a summary table, and a leaf node also contains a I_m table with $(f+1)^2$ entries and a heap in size f for fast locating and merging the entries. I_m table keeps the value $I_m(i,j)$ for any pair of entries. Together with the heap, it is fast to keep track of the minimum I_m. An internal node (non-leaf) in the tree contains only the aggregation information of its child nodes – each entry in the internal node points to a child node and its summary table is the sum of the summary tables in the child node.

Let a summary table represented with a vector $\vec{s}$ and the entropy characteristic of any internal node C_i denoted as $EC_i(n_i, \vec{s_i})$, where n_i is the number of records summarized by this node. Let C_{ij}, $1 \leqslant j \leqslant f$ represent the child nodes of C_i. HE-Tree maintains the following property.

$$EC_i(n, \vec{s}) = \sum_{j=1}^{f} EC_{ij}(n_{ij}, \vec{s_{ij}}) = EC_i(\sum_{j=1}^{f} n_{ij}, \sum_{j=1}^{f} \vec{s_{ij}})$$

i.e., the parent node represents the mergence of the child nodes. The key of HE-Tree is to approximately minimize the overall expected entropy by locally minimizing the expected entropy of the selected branch $\bar{H}(C_i^f)$ in each insertion of new record. This local minimization is achieved through the following algorithms in constructing the HE-Tree.

4.2 Constructing HE-Tree The construction of HE-Tree consists of two phases: *the growing phase* and *the absorbing phase*. The algorithms for construction are carefully designed to minimize the expected entropy of the subclusters and adapt to the change of entropy in data stream.

Growing Phase. In the growing phase, the tree grows until the number of leaf nodes reaches $\lceil n_c/f \rceil$. When a new coming record is inserted into the existing tree, the first subroutine is for locating the target leaf node for insertion/absorption. Let e denote the inserted record and e_i denote one of the entry in current node. The search begins at the root node. Since each entry in the internal node is the summarization of its sub-tree, we can find the most similar entry to e by finding the minimum I_m among $I_m(e, e_i), i = 1..f$, i.e.

$$e_t = argmin_{e_i}\{I_m(e, e_i), i = 1..f\}$$

Iteratively, the same criterion is applied to the selected child node until a leaf node is reached. If the target leaf node has empty entries and $I_m(e, e_i) \neq 0$ for all occupied entries, the record occupies one empty entry. Otherwise, the new record is merged to the identical entry. We give the sketch of the subroutines in Algorithms 1 and 2.

When the target leaf node is full, a split operation is applied. In split algorithm, we partition the entries into two groups. First, two pivot entries (e_r, e_s) is found in

Algorithm 1 HE-Tree.locate(node, e)

```
node ← target node, e ← target entry
if node is leaf then
  return node
end if
for Each entry e_i in node do
  I_m^i ← I_m(e, e_i)
end for
e_t ← argmin_{e_i}{I_m^i}
return locate(e_t.subtree, e)
```

Algorithm 2 HE-Tree.insert(node, e)

```
e ← inserted entry, node ← target node
for Each entry e_i in node do
  if I_m(e, e_i) == 0 then
    merge(e, e_i), return
  end if
end for
if node.have_empty_entry() then
  node.enter(e)
  if (node.num_entry() == f − 1 and (not tree_full()
  or is_internal(node)) then
    split(node)
  end if
else
  leaf-merging(node, e) //fine merging in absorption
  phase
end if
```

Algorithm 3 HE-Tree.split(node)

```
node ← target node
(e_a, e_b) ← argmax_{(e_i,e_j)}{I_m(e_i, e_j)}
partition_a ← e_a, partition_b ← e_b
for Each entry e_i in node do
  if I_m(partition_a, e_i) < I_m(partition_b, e_i) then
    partition_a ← partition_a ∪ e_i
  else
    partition_b ← partition_b ∪ e_i
  end if
end for
if is_leaf(node) and not done then
  re-insert(root, entries in partition_a)
else
  newnode ← partition_a, remove(node, partition_a)
  e_new ← summary(newnode), insert(node.parent,
  e_new)
end if
```

the target node that have the maximum I_m if merging them – they are regarded as the most dissimilar pair among all pairs.

$$(e_r, e_s) = argmax_{e_r,e_s}\{I_m(e_r, e_s), i = 1..f\}$$

The two pivot entries then become the two seed clusters. The rest entries are sequentially assigned to the two clusters so that the overall expected-entropy of the partition keeps minimized. A new node is generated accommodating one of the two sets of entries. At the same time, one entry is added into the parent node pointing to the new node. The insertion/splitting continues until the number of leaf entries reaches n_c. Algorithm 3 gives the detailed description.

Absorbing Phase. In the second phase, the same locating algorithm is applied to locate the target leaf node for the new record. However, we have no insertion allowed since the entries are all occupied. Instead, in the leaf node we need to merge the most similar two items among the f+1 items (micro-merging) – the f entries in the leaf node plus the new record. This allows the tree to rapidly adapt to the change of clustering structure in the entry level. In each leaf node, we maintain a I_m table and a heap for the f entries. When a new record comes, only f calculations of incremental entropy are needed to update the I_m table and the heap, before selecting the most similar two to merge.

The locating algorithm with the Incremental Entropy criterion will assign a new record to the approximately best leaf node and a fine micro-merging will well adapt the local structural change happening within the node. Experiments show that the summary entries at the leaf nodes together can precisely describe the global clustering structure.

4.3 Setting of Parameters The setting of the two parameters f and n_c can affect the efficiency and quality of summarization. Let h be the height of the tree (root is at level 1). For simplicity, we always construct full trees and allow $n_c = f^h$ to vary from hundreds to thousands. For example, for $f = 15$, we can either use a two-layer tree, where the number of leaf entries $n_c = 225$, or a three-layer tree where $n_c = 3375$. In experiment, we show that a smaller f always results in faster summarization, but can undermine the quality of summarization when the clustering structure is changing. A small f may cause more imprecise mergence to happen in the second phase, since the less entries the lower level of precision is guaranteed for absorption. Larger f with the same height of tree will increase the cost – $O(dmf)$ in absorbing phase. On the other hand, larger f increases the ability adapting to the change of clustering structure since we can do more precise merging in the absorbing phase. To tradeoff the performance and robustness, we can set the tree to be

2 ∼3 layers, with $f = 10 \sim 20$.

4.4 Complexity of HE-Tree The time complexity of constructing HE-Tree can be divided into two phases. In the growing phase, about f^h records are inserted into the tree and each record needs at most $O(hf)$ comparison to locate the target node. In the absorption phase, besides the cost of locating, each record needs micro-merging at leaf, which costs $O(f)$ incremental-entropy calculation. Incremental-entropy involves only weighted entropy which costs O(dm). Therefore, the cost is $O((h + dm)f)$. Since f is usually a small value, e.g. $10 \sim 20$ and $h = 2$ or 3 in practice. Thus, the total cost is only dominated by the number of dimensions d and the average cardinality m of the dataset, i.e. the factor dm.

There are $O(f^h)$ nodes in the tree. Each leaf node needs approximately $O(fdm + f^2)$ space, where the summary table for each entry needs $O(dm)$ and the I_m table needs $O(f^2)$. Each internal node needs only $O(fdm)$, holding the summary tables and the pointers to the children nodes. Approximately, a HE-Tree needs $O((dm+f)f^{h+1})$ space. With fixed small f and h, again only the factor dm of the dataset determines the size of the tree. Except the datasets having very large dm, e.g. over 10k, HE-Tree usually needs small amount of memory.

In summary, HE-Tree can efficiently summarize the entropy characteristics of the data stream with small amount of time and space cost.

5 Framework for Monitoring the Change of Clustering Structure

In the last section, we have designed the HE-Tree algorithm for summarizing the data stream. In this section, we first briefly describe how to take the summary of data stream and generate an effective BkPlot. Then, we present the entire framework for detecting the change of clustering structure.

Extended ACE Algorithm The extended ACE algorithm is a hierarchical clustering algorithm built on the base clusters generated by the HE-Tree, i.e., the summary information in the leaf entries of the HE-Tree. Suppose there are n_c sub-clusters generated by the summarization. It consecutively merges the pair of clusters that minimizes the Incremental Entropy among the remaining clusters. With the help of similar structures of summary table, I_m table and heap, as used in the absorption phase of HE-Tree construction, the extended ACE algorithm can be optimized to have $O(n_c^2)$ Incremental Entropy calculation. Since n_c is only several hundreds in practice, the extended ACE algorithm can be done very quickly. We have shown that the basic ACE algorithm [11] can generate high-quality BkPlots, and we will show that the extended ACE algorithm with the HE-Tree summarization can also effectively identify the best Ks for categorical data stream.

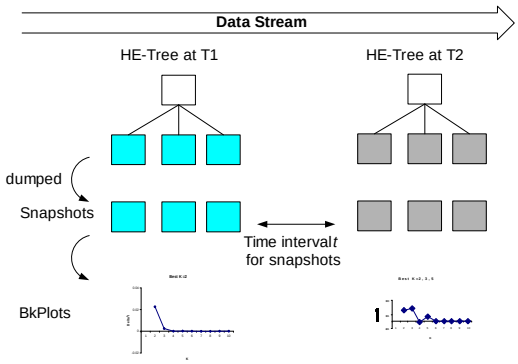

Figure 5: Framework for detecting change of clustering structure in categorical data streams.

A Brief Framework With the HE-Tree and the extended ACE algorithm, we can precisely monitor the change of clustering structure in the categorical data stream. The framework is illustrated in Figure 5. The working mechanism can be described as follows.

1. The records from the data stream are inserted into the HE-Tree by order. Each insertion costs $O((h + dm)f)$;
2. At certain time interval Δt, the summary tables in the leaf nodes are dumped out (to a piece of memory or to hard disk). It only costs $O(dmn_c)$ bytes to store each of such snapshots;
3. the extended ACE algorithm are performed on the snapshot as soon as it is dumped, the result of which generates a BkPlot. The cost is $O(dmn_c^2)$ [11].

Basically, the cost of first step restricts how many records the framework can monitor in unit time. As we have shown, f also affects the precision of summarization. There is a tradeoff between the precision and the capacity of monitoring system, tuned by the parameter f. The cost of third step affects how frequently we can generate a BkPlot. The time interval Δt in the second step is directly determined by the cost of generating BkPlot, i.e., Δt should be greater than $O(dmn_c^2)$. Reducing n_c allows more snapshots to be processed in unit time, and thus more details about the changes to be observed. In practice, $n_c = 400 \sim 1000$ is enough

to generate precise BkPlot, which means $O(dmn_c^2)$ is usually a small number. How often we need to monitor the change of stream also depends on the application requirement. For example, in monitoring communication network, we may need to detect the changes between seconds. However, for real-world traffic monitoring, we may only need to check the change in every ten minutes.

The neighboring BkPlots are analyzed to see the difference between the clustering structure. BkPlots can be represented as a function $B(K)$, where K is the number of clusters and the distinctive $B(K)$s indicate the candidate best Ks. Without loss of generality, we suppose the first κ distinctive Ks on BkPlots are $\Gamma = \{k_1, k_2, \ldots k_\kappa\}$. Let Γ^{old} and Γ^{new} represent two set of Ks on the consecutive BkPlots, respectively. There are two kinds of important differences we need to notice.

1. If Γ^{old} and Γ^{new} are not identical, the clustering structure is dramatically changed, which raises an "alarm" that we need to analyze the snapshot of Γ^{new} in detail.

2. If Γ^{old} and Γ^{new} are identical, but at certain k_i that $|B(k_i^{new}) - B(k_i^{old})| > \theta$, where θ is a threshold we need to notice, we can infer some minor changes in clustering structure – if $B(k_i^{new}) > B(k_i^{old})$, i.e., ECG curve changes more dramatic at the critical k_i, the clusters grows distinctively; otherwise, the boundaries between clusters become vague and some clusters tend to converge.

6 Experimental Results

The goal of the experiments is two-fold. 1) We investigate the parameter setting of HE-Tree and give the estimate of appropriate settings; 2) We want to show that HE-Tree summarization together with the extended ACE algorithm can provide high-quality monitoring result.

Datasets We construct a synthetic dataset DS1 with the following way, so that the clustering structure can be intuitively identified and manually labeled before running the experiments. The synthetic dataset has a two-layered clustering structure (Figure 6) with 30 attributes and N rows. It has four same-sized clusters in the top layer. Each cluster has random categorical values selected from {'0','1','2','3','4', '5'} in some distinct set of attributes (the dark area in Figure 6), while the rest attributes are set to '0'. Two of the four clusters also have clustering structure in the second layer. This synthetic data has clearly defined clustering structure, and each record in the dataset distinctly belongs to one cluster. This dataset is primarily used in exploring the effect of the parameters of HE-Tree to the precision of clustering result and the efficiency of summarization. And later, it is also used to demonstrate a sequence of monitoring results. We also use a real dataset: "US

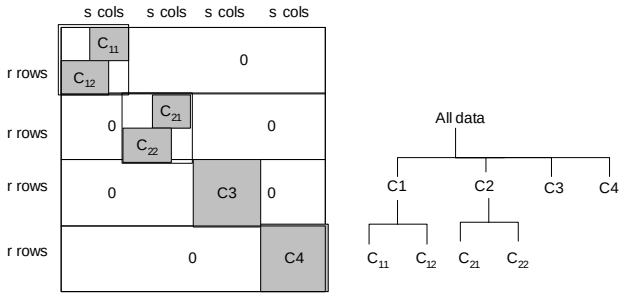

Figure 6: Clustering structure of DS1

Census 1990 Data " in the experiment. This dataset is a discretized version of the raw census data, originally used by [26]. It can be found in UCI KDD Archive [1]. Many of the less useful attributes in the original data set have been dropped, the few continuous variables have been discretized and the few discrete variables that have a large number of possible values have been collapsed to have fewer possible values. The total number of preserved attributes is 68. The first attribute is the sequence number, which is discarded in clustering. This dataset is very large, containing about 2 million records. By visualizing its sample dataset, we can observe three clusters, two of which are close to each other (Figure 7.

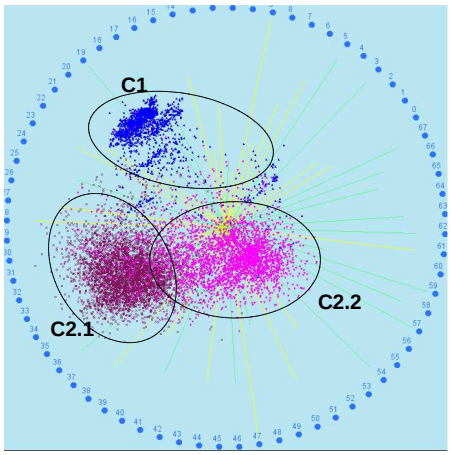

Figure 7: Clustering structure of DS1

Error Rate The cluster labels in the synthetic dataset DS1 allow us to evaluate the quality of clustering result more accurately by using the *Error Rate* measure. Suppose the best K is identified. The Error

[1] http://kdd.ics.uci.edu/

Rate is defined based on the *confusion matrix*, where each element c_{ij} $1 \leqslant i, j \leqslant K$ represents the number of points from the labeled cluster j assigned to cluster i by the algorithm. Let {(1), (2), ..., (K)} be any permutation of sequence {1, 2, ..., K}. There is a permutation that maximizes the number of consistent points m_c, which is defined as follows.

$$m_c = \max\{\sum_{i=1}^{K} c_{i(i)}, \text{for any } \{(1), (2), \ldots, (K)\}\}$$

We define Error Rate as $1 - \frac{m_c}{N}$, N is the total number of points.

6.1 Parameter Setting for HE-Tree For a full tree, the fan-out f of tree node and the height of the tree determine the tree structure. To simplify the investigation and maximize the quality of the summarization, we always use full trees in the experiment. Intuitively, for a fixed f, the higher tree (the larger h), the finer granularity of summarization will be delivered. However, most likely we care about only the clustering structures having less than 20 clusters. Therefore, a short tree, which generates less than one thousand sub-clusters, is enough for achieving high-quality BkPlot with ACE clustering algorithm. The experiments will be focused on the full short trees (e.g., $h = 2$) with varying fan-out f from 10 to 30. A set of datasets (20 datasets) in the same structure shown in Figure 6 are generated, the result is statistically based on the 20 runs.

Figure 8 shows the cost for HE-Tree summarization is linear and the cost also varies linearly according to f, which is consistent with our analysis. Figure 9 shows the effect of different summarization structures to the quality of final clustering result for "Unordered DS1". Unordered DS1 randomly stores the records from different clusters,i.e., there is little change of clustering structure in processing the entire data stream. The result shows some variances between the error rates for different f, but overall the error rates are similar and low.

"Ordered DS1" shows a more interesting scenario, where clustering structure dramatically changes when time goes by. In such situations, f may significantly affect the quality of monitoring. Figure 10 shows the result of sequentially processing the clusters C_{11} to C_4. A tree with larger f seems more adaptive to the change of clustering structure. The reason that HE-Tree is more sensitive to the setting of f when the clustering structure changes dramatically can be understood as follows. The initial records from the same cluster already occupy the slots in the growing stage. When a new cluster emerges, since there is no entry in the tree belonging to the new cluster, new slots are given to the new records by merging other similar entries, or, some initial records may be absorbed to other clusters by small chance, which causes the error. It shows that increasing f from 10 to 20 can considerably reduce the error, but $f = 30$ will not significantly improve the result of $f = 20$. Balanced with the time cost and the robustness, $f = 20$ seems the best for efficiently adapting the change of structure.

6.2 Robustness of BkPlots by HE-Tree/Extended ACE In this experiment, we want to compare the accuracy of BkPlots generated by ACE on small sample set and by HE-Tree/extended ACE on large stream data. We run the experiment on both the synthetic data and the real US Census data. The small sample size is set to 500 for ACE, and large samples sizes are 10K and 100K. The sample sets are uniformly drawn from the original dataset, therefore, they are supposed to have the same clustering structure.

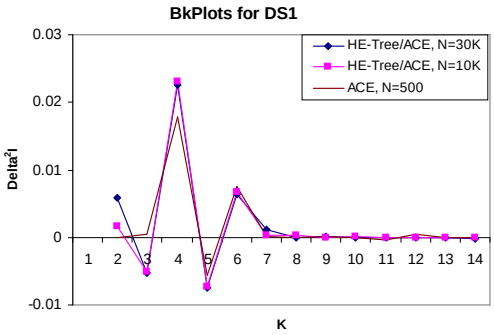

Figure 11: BkPlots for DS1.

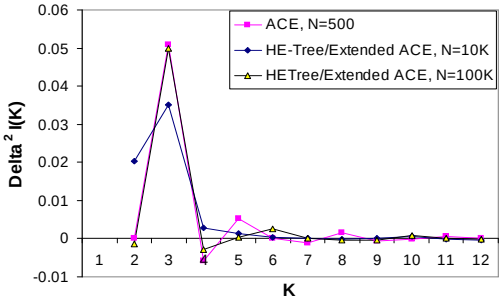

Figure 12: BkPlots for Census.

Figure 11 for DS1 shows all of the three BkPlots can identify the primary best Ks: 4 and 6, while a little noise appears at K=2 when the sample size is large. All BkPlots for Census data (Figure 12) strongly suggest the best K=3, while K=2 is probably another candidate(which groups the cluster C2.1 and C2.2 together).

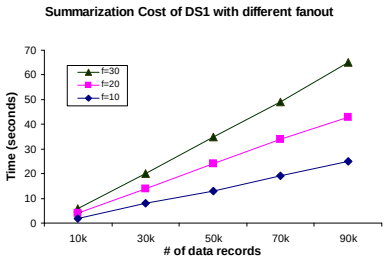

Figure 8: Cost of HE-Tree summarization with different fanout f.

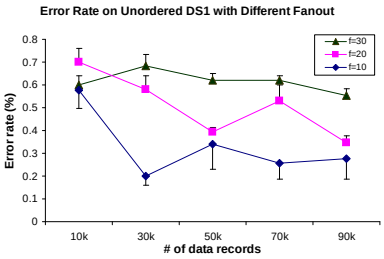

Figure 9: Error rate of ACE clustering result with HE-Tree summarization on randomly ordered records

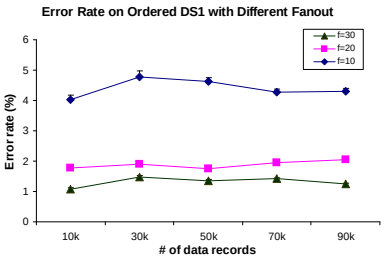

Figure 10: Error rate of ACE clustering result with HE-Tree summarization on ordered records

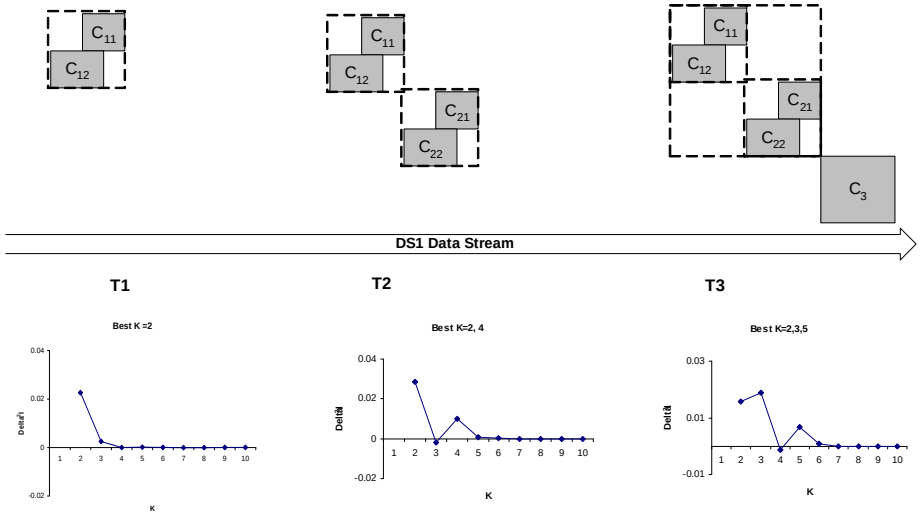

Figure 13: Monitoring DS1-stream.

The result confirms that HE-Tree summarization can preserve the primary clustering structure and thus HE-Tree combined with ACE method is a robust approach for monitoring the change of clustering structure.

6.3 Monitoring the Changes We also demonstrate the progressive monitoring results of the two data streams: DS1-stream and Census-stream . The DS1-stream simulates the 4/6-cluster structure shown in Figure 6. The clusters enter the stream in the sequence of C_{11}, C_{12},C_{21},C_{22}, C_3, and C_4. Each of the small clusters have 5K records and each of the large clusters have 10K records. Snapshots are saved at N=10K, 20K, and 30K, respectively.

The progressive results for DS1-stream in Figure 13 clearly identify the change of clustering structure. At T1:N=10K, C_{11} and C_{12} have been present at the stream, thus two clusters are identified. At T2:N=20K, C_{21} and C_{22} emerge and the two-layer structure is identified (the best K = 2, 4). At T3:N=30K, C_3 appears, and the BkPlot detects that the primary two-layer structure is changed to K=3, 5, while the BkPlot also suggest an additional layer at K=2, which consists two cluster ($C_{11}, C_{12}, C_{21} and C_{22}$) and ($C_3$).

We partition the census dataset into four parts and mix the parts sequentially so that the special clustering structures appear in different stage as Figure 14 shows. At first snapshot, there are clearly two clusters; in the second one, the third cluster shows vaguely; finally, a two-layer clustering structure (K=2 and 3) appears in the third snapshot. Snapshot 2 demonstrates that the HE-Tree can also capture the fine changes in clustering structure very well.

7 Related Work

While many numerical clustering algorithms [23, 24] have been published, only a handful of categorical

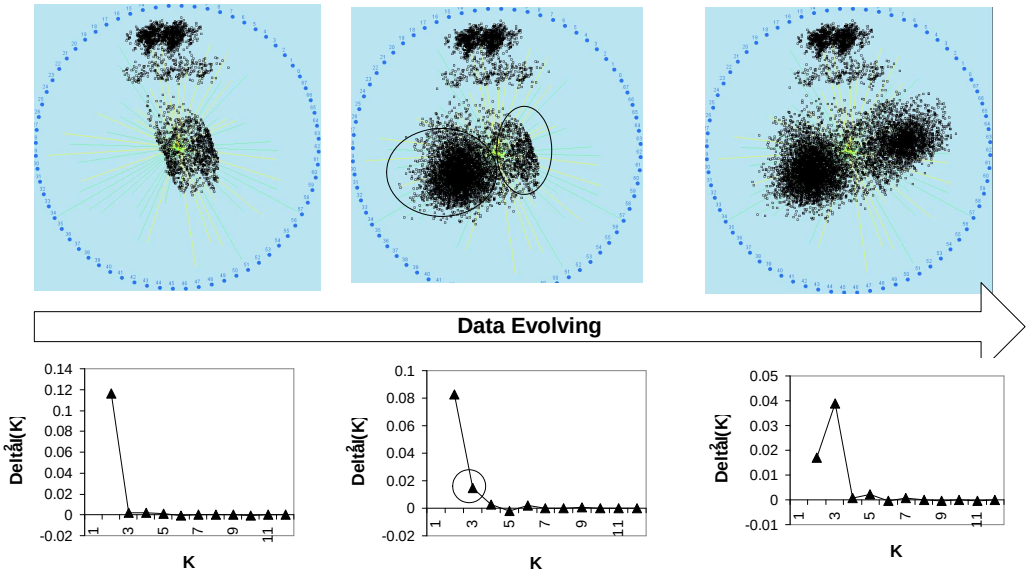

Figure 14: Monitoring Census-stream.

clustering algorithms appear in literature. Although it is unnatural to define a distance function between categorical data or to use the statistical center (the mean) of a group of categorical items, there are some algorithms, for example, K-Modes [22] algorithm and ROCK [20] algorithm, trying to fit the traditional clustering methods into categorical data. However, since the numerical similarity/distance function may not describe the categorical properties properly and intuitively, it leaves little confidence to the clustering result. CACTUS [15] also partly adopts the linkage idea used in ROCK.

Gibson et al. introduced STIRR [18], an iterative algorithm based on non-linear dynamical systems. STIRR represents each attribute value as a weighted vertex in a graph. Starting with the initial conditions, the system is iterated until a "fixed point" is reached. When the fixed point is reached, the weights in one or more of the "basins" isolate two groups of attribute values on each attribute. Even though they proved this approach works for some experimental datasets having two partitions, the user may hesitate in using it due to the complicated and not intuitive working mechanism.

Coolcat [4] is kind of similar to KModes. However, Coolcat assigns the item to the cluster that minimizes the expected entropy. Considering the cluster centers may shift, a number of worst-fitted points will be re-clustered after a batch. Li et al [25] proposed a Monte-Carlo method to minimize the expected entropy, which is slower than Coolcat but can be more likely to achieve the sub-optimal results. Cross Association [9] tries using MDL to partition boolean matrix along row direction and column direction at the same time. In fact, MDL is equivalent to entropy criterion as [25] shows. Some closely related work also borrows concepts from information theory, including Co-clustering [14], Information Bottleneck [27] and LIMBO [3]. The results all show that the entropy related concepts work very well for categorical data.

Clustering data streams becomes one of the important technique for analyzing the data streams [19]. In [2], a framework CluStream is proposed for clustering evolving *numerical* data streams, which mainly concerns summarizing and storing the sketch of the data stream. There has been recent work on frameworks based on velocity density estimation [1]. Clustering categorical data stream was first addressed by Coolcat[4], but no more related issues such as detecting the change of clustering structure are addressed yet. Nonparametric testing is used in paper [5], to detect the changes in data streams, and we propose that monitoring the change of clustering structure is also very useful. There is also other work using various statistical testing [16, 17] to monitor the changes in data stream, and change detection in semi-structured data [10].

8 Conclusion

In this paper, we address the problem of detecting the change of clustering structure in categorical data streams with a novel framework. The key of the framework is the combination of BkPlot method and Hierarchical Entropy Tree (HE-Tree) summarization

structure and algorithms. HE-Tree is designed as a memory-efficient structure − the tree is usually a short tree (height = 2 or 3) with small number of leaf nodes, which store the information of the summarized sub-clusters. In order to observe the change of clustering structure, snapshots of the leaf entries of HE-Tree are dumped in certain time interval, which is then processed by the extended ACE clustering algorithm to generate high-quality BkPlots, with which we can easily identify whether and how the clustering structure in the stream is changed. Experiments show with HE-Tree and BkPlot method we can effectively detect the change of critical clustering structure in categorical data streams.

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